

Constructing partial geometries with prescribed automorphism groups^{*}

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Introduction

A **partial geometry** $pg(s, t, \alpha)$ is an incidence structure $(\mathcal{P}, \mathcal{L}, I)$ such that:

- every line is incident with $s + 1$ points,
- every point is incident with $t + 1$ lines,
- every pair of points is incident with at most one line,
- for every non-incident point-line pair (P, ℓ) , there are exactly α points on ℓ collinear with P .

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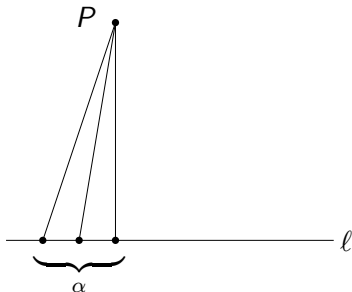
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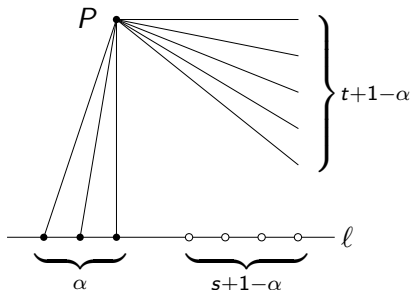
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$$\alpha \leq \min\{s + 1, t + 1\}$$

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The **dual** of a $pg(s, t, \alpha)$ is a $pg(t, s, \alpha)$. In the sequel we will always assume $\alpha < s \leq t$.

Necessary existence conditions

Counting arguments in a $pg(s, t, \alpha)$ yield:

$$v := |\mathcal{P}| = (s + 1) \frac{(st + \alpha)}{\alpha}, \quad b := |\mathcal{L}| = (t + 1) \frac{(st + \alpha)}{\alpha}.$$

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The **point graph** of a $pg(s, t, \alpha)$ is strongly regular with parameters

$$SRG(v, s(t + 1), s - 1 + t(\alpha - 1), \alpha(t + 1)).$$

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From this we get necessary conditions on the parameters:

$$\alpha \mid (s + 1)st \quad \text{and} \quad (t + 1)st \quad (\text{integrality of } v \text{ and } b)$$

$$\alpha(s + t + 1 - \alpha) \mid st(s + 1)(t + 1) \quad (\text{multiplicities of eigenvalues of SRG})$$

$$(s + 1 - 2\alpha)t \leq (s - 1)(s + 1 - \alpha)^2 \quad (\text{Krein inequalities for SRG})$$

Admissible parameters of small $pg(s, t, \alpha)$ ($v \leq 100$)

Part.geom.	Npg	Point gr.	Nsrg	Line gr.	Nsrg
$pg(2, 2, 1)$	1	(15, 6, 1, 3)	1	(15, 6, 1, 3)	1
$pg(2, 4, 1)$	1	(27, 10, 1, 5)	1	(45, 12, 3, 3)	78
$pg(3, 4, 2)$	0	(28, 15, 6, 10)	4	(35, 16, 6, 8)	3854
$pg(3, 3, 1)$	2	(40, 12, 2, 4)	28	(40, 12, 2, 4)	28
$pg(4, 6, 3)$	2	(45, 28, 15, 21)	1	(63, 30, 13, 15)	+
$pg(3, 5, 1)$	1	(64, 18, 2, 6)	167	(96, 20, 4, 4)	+
$pg(5, 8, 4)$	0	(66, 45, 28, 36)	1	(99, 48, 22, 24)	+
$pg(6, 6, 4)$?	(70, 42, 23, 28)	+	(70, 42, 23, 28)	+
$pg(4, 7, 2)$	0	(75, 32, 10, 16)	0	(120, 35, 10, 10)	?
$pg(3, 6, 1)$	0	(76, 21, 2, 7)	0	(133, 24, 5, 4)	?
$pg(5, 5, 2)$	≥ 1	(81, 30, 9, 12)	+	(81, 30, 9, 12)	+
$pg(4, 4, 1)$	1	(85, 20, 3, 5)	+	(85, 20, 3, 5)	+
$pg(6, 10, 5)$?	(91, 66, 45, 55)	1	(143, 70, 33, 35)	+
$pg(4, 9, 2)$	0	(95, 40, 12, 20)	0	(190, 45, 12, 10)	?
$pg(5, 6, 2)$?	(96, 35, 10, 14)	?	(112, 36, 10, 12)	?
$pg(5, 9, 3)$	0	(96, 50, 22, 30)	0	(160, 54, 18, 18)	?

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Partial geometries with prescribed automorphism groups

An **automorphism** of a partial geometry \mathcal{G} is a permutation of the points mapping lines onto lines. The **full automorphism group** is denoted $\text{Aut}(\mathcal{G})$. Any subgroup $G \leq \text{Aut}(\mathcal{G})$ is called an **automorphism group** of \mathcal{G} .

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Construction method:

- 1 Choose a permutation group G on the set of points \mathcal{P} .
- 2 Compute the orbits of G on k -element subsets of \mathcal{P} .
- 3 Select orbits comprising lines of the partial geometry.

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B. McKay, A. Piperno, *nauty and Traces*, <https://pallini.di.uniroma1.it>

GAP – Groups, Algorithms, and Programming, <http://www.gap-system.org>

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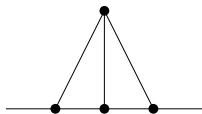
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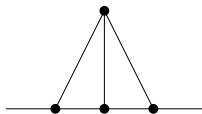
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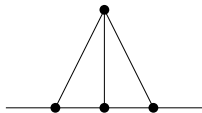
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How to generate the orbits?

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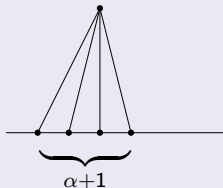
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Theorem.

A partial linear space of order (s, t) is a $pg(s, t, \alpha)$ if and only if it does not contain the configuration:



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We want to select a subset of the orbits $\mathcal{O}_1, \dots, \mathcal{O}_n$ such that $\bigcup_{i=1}^n \mathcal{O}_i$ is the set of lines of a $pg(5, 5, 2)$. Necessary: $\sum_{i=1}^n |\mathcal{O}_i| = 81$.

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Define the **compatibility graph** with the 181 orbits as vertices and their sizes as weights. Orbits \mathcal{O}_1 and \mathcal{O}_2 are joined by an edge if $|X \cap Y| \leq 1$ for all $X \in \mathcal{O}_1$, $Y \in \mathcal{O}_2$ and $\mathcal{O}_1 \cup \mathcal{O}_2$ does not contain the forbidden configuration.

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Search for cliques of weight $b = 81$ in the compatibility graph.

S. Niskanen, P. R. J. Östergård, *Cliquer user's guide, version 1.0*, Communications Laboratory, Helsinki University of Technology, Espoo, Finland, Tech. Rep. T48, 2003.

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The compatibility graph has 181 vertices and 528 edges (density 3.2%).

There are 384 cliques of weight 81, all of them correspond to $pg(5, 5, 2)$.

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The compatibility graph has 181 vertices and 528 edges (density 3.2%).

There are 384 cliques of weight 81, all of them correspond to $pg(5, 5, 2)$.

Using `nauty`, one finds that there are **two** non-isomorphic $pg(5, 5, 2)$.

$|\text{Aut}(\mathcal{G}_1)| = 58\,320$, isomorphic to the geometry of van Lint and Schrijver.

$|\text{Aut}(\mathcal{G}_2)| = 972$, a new partial geometry $pg(5, 5, 2)$!

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V. Krčadinac, *A new partial geometry $pg(5, 5, 2)$* , 16 September 2020.

<https://arxiv.org/abs/2009.07946>

We give a computer-free description of the new $pg(5, 5, 2)$ using a four-dimensional vector space over $GF(3)$, by changing some lines of the geometry of van Lint and Schrijver.

A general construction of partial geometries

Let \mathcal{R} be a projective plane of order q . A subset \mathcal{A} of points is called a d -arc if d is the greatest number of collinear points in \mathcal{A} .

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The number of points in a d -arc is $|\mathcal{A}| \leq dq - q + d$. Equality holds if and only if every line is either disjoint from \mathcal{A} , or intersects \mathcal{A} in exactly d points. In this case \mathcal{A} is called a maximal arc (for $d = 2$ a hyperoval).

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Theorem (J. A. Thas, W. D. Wallis, 1973).

Let \mathcal{A} be a maximal d -arc in a projective plane of order q . The set of d -secants of \mathcal{A} as POINTS and the set of points not in \mathcal{A} as LINES constitute a partial geometry $pg(s, t, \alpha)$ for $s = q(d - 1)/d$, $t = q - d$, and $\alpha = (q - d)(d - 1)/d$.

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We want $s \leq t$ and hence only consider $d \leq \sqrt{q}$.

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R. Mathon, *The partial geometries $pg(5, 7, 3)$* , *Congr. Numer.* **31** (1981), 129–139.

Mathon proved that there are precisely **two** $pg(4, 6, 3)$ up to isomorphism. The one above has $\text{Aut}(\mathcal{G}) = PGL(2, 8) \rtimes \mathbb{Z}_3$ of order 1512. The other one does not come from a hyperoval in $PG(2, 8)$ and has $\text{Aut}(\mathcal{G}) = (\mathbb{Z}_3^2 \rtimes Q_8) \rtimes \mathbb{Z}_3$ of order 216. Both geometries can be obtained from a group of order 18 isomorphic to $S_3 \times \mathbb{Z}_3$. This is the “largest common subgroup” of their full automorphism groups.

A general construction of partial geometries

- $q = 9$ There are **4** projective planes: $PG(2, 9)$, Hall, dual Hall and Hughes.
- $d = 3$ The four planes do not contain maximal 3-arcs ($\rightsquigarrow pg(6, 6, 4)$).

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$q = 16$ There are 22 known projective planes.

$d = 2$ Hyperovals exist in 18 known planes $\rightsquigarrow pg(8, 14, 7)$ (93 non-isom.)

T. Penttila, G. F. Royle, M. K. Simpson, *Hyperovals in the known projective planes of order 16*, J. Combin. Des. **4** (1996), no. 1, 59–65.

M. Gezek, V. D. Tonchev, *On partial geometries arising from maximal arcs*, 30 August 2020. <https://arxiv.org/abs/2008.13246>

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$ \text{Aut} $	$\#pg(8, 14, 7)$	$ \text{Aut} $	$\#pg(8, 14, 7)$	$ \text{Aut} $	$\#pg(8, 14, 7)$
16320	1	64	8	6	2
320	1	32	3	4	4
144	1	16	59	3	1
112	2	14	2	2	3
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$q = 16$ There are 22 known projective planes.

$d = 4$ Maximal 4-arcs exist in 18 known planes (not completely classified)
 $\rightsquigarrow pg(12, 12, 9)$ (59 non-isomorphic)

A general construction of partial geometries

$ \text{Aut} $	$\#pg(12, 12, 9)$	$ \text{Aut} $	$\#pg(12, 12, 9)$
408	1	32	21
144	3	24	12
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68	1	4	5
48	4		

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Work in progress...



- find $pg(6, 6, 4)$ or prove that they do not exist
- find $pg(8, 14, 7)$ not arising from a hyperoval
- find more $pg(12, 12, 9)$
- ...

Thanks for your attention!